## **Chapter 2: Memory Management: Early Systems**

## TRUE/FALSE

1.	Early memory management schemes are still used in today's operating systems.						
	ANS: F	PTS:	1	REF:	32		
2.	. All computers have only a finite amount of memory and if a program doesn't fit, then either the size the main memory must be increased or the program must be modified.						
	ANS: T	PTS:	1	REF:	32		
3.	3. To overlay is to transfer segments of a program from main memory into secondary storage for execution, so that two or more segments take turns occupying the same memory locations.						
	ANS: F	PTS:	1	REF:	32		
4. The first step in loading a job in a single-user system is storing the first memory location of printo the base register (for memory protection).							
	ANS: T	PTS:	1	REF:	33		
5.	A single-user system	suppor	ts multiprogran	nming.			
	ANS: F	PTS:	1	REF:	33		
6.	The first attempt to a	llow for	r multiprogram	ming u	sed fixed partitions.		
	ANS: T	PTS:	1	REF:	34		
7. The problem of partition intrusion is present in single-user contiguous allocation schemes.							
	ANS: F	PTS:	1	REF:	34		
8. The algorithm used to store jobs into memory requires a few more steps than the one used for single-user system because the size of the job must be matched with the size of the partition to sure it fits completely.							
	ANS: T	PTS:	1	REF:	34		
9.	The fixed partition scheme does not require that the entire program be stored contiguously and in memory from the beginning to the end of its execution.						
	ANS: F	PTS:	1	REF:	35		
10. The fixed partition scheme works well if all of the jobs run on the system are of the same sizes are known ahead of time and don't vary between reconfigurations.							
	ANS: T	PTS:	1	REF:	35		

11.	In a fixed partition scheme, large jobs may have a longer turnaround time as they wait for free partitions of sufficient size or may never run.						
	ANS: T	PTS:	1	REF:	36		
12.	2. The best-fit allocation method keeps the free/busy lists organized by memory locations, low-omemory to high-order memory.						
	ANS: F	PTS:	1	REF:	38		
13.	A large job can hav	e problei	ns with a first-	fit mem	nory allocation list.		
	ANS: T	PTS:	1	REF:	39		
14.	The best-fit free list to implement.	scheme	uses memory n	nore ef	ficiently than the first-fit free scheme but it is slower		
	ANS: T	PTS:	1	REF:	40		
15.	The first-fit algorith blocks.	nm assum	nes that the Men	mory M	Ianager keeps only one list containing free memory		
	ANS: F	PTS:	1	REF:	40		
16.	One of the problems with the best-fit algorithm is that the entire table must be searched before the allocation can be made because the memory blocks are physically stored in sequence according to their location in memory.						
	ANS: T	PTS:	1	REF:	42		
17.	7. Research continues to focus on finding the optimum allocation scheme.						
	ANS: T	PTS:	1	REF:	43		
18.	18. For a fixed partition system, memory deallocation is quite complex.						
	ANS: F	PTS:	1	REF:	44		
19.	9. A null entry in the busy list occurs when a memory block between two other busy memory bl returned to the free list.						
	ANS: T	PTS:	1	REF:	47		
20.	In the relocatable dynamic partitions scheme, the Memory Manager relocates programs to gather together all of the empty blocks and compact them to make one block of memory large enough to accommodate some or all of the jobs waiting to get in.						
	ANS: T	PTS:	1	REF:	48		
21.	Memory is allocated	d during	garbage collect	ion.			
	ANS: F	PTS:	1	REF:	48 49		

22. During compaction, the operating system must distinguish between addresses and data validations are not obvious once the program has been loaded into memory.							
	ANS: T	PTS: 1	REF:	49			
23.	After relocation and compaction, both the free list and the busy list are updated.						
	ANS: T	PTS: 1	REF:	51			
24.	The bounds register is used to store the highest (or lowest, depending on the specific system) location in memory accessible by each program.						
	ANS: T	PTS: 1	REF:	52			
25.	Compaction should always be performed only when there are jobs waiting to get in.						
	ANS: F	PTS: 1	REF:	53 54			
MUL	TIPLE CHOICE						
1.	<ul> <li>The following,, describes the first memory allocation scheme.</li> <li>a. Each program to be processed was loaded into secondary storage, then swapped into memory in parts</li> <li>b. Each program to be processed was partially loaded into memory, then granted more memory as needed</li> <li>c. Each program to be processed was allocated a portion of memory and could negotiate with other programs to access more memory</li> <li>d. Each program to be processed was loaded in its entirety into memory and allocated as much contiguous space in memory as it needed</li> </ul>						
	ANS: D	PTS: 1	REF:	32			
2.	In a single-user sy a. sequentially b. intermittently	stem, jobs are prod	c.	randomly in order of longest job to shortest job			
	ANS: A	PTS: 1	REF:	32			
<ul> <li>3. In the algorithm to load a job in a single-user system, the program counter is initially s</li> <li>a. the address of the last memory location</li> <li>b. the number of instructions</li> <li>c. zero</li> <li>d. the address of the first memory location</li> </ul>							
	ANS: D	PTS: 1	REF:	33			
4.	Fixed partitions ar a. complete b. static	re also called	c.	direct sized			
	ANS: B	PTS: 1	REF:	34			
5.	a. Comparing job			in a fixed partition.			

	c.	Setting counter t	Determining the job's requested memory size etting counter to one Placing the job in a waiting queue					
	AN	IS: B	PTS:	1	REF:	34		
6.	<ul> <li>6. In the partition scheme, the table that the Memory Manager uses to keep track of jobs is compthe</li> <li>a. partition size, memory address, and status</li> <li>b. status, access, and memory address</li> <li>c. partition size, status, and access</li> <li>d. partition size, memory address, access, and status</li> </ul>							
	AN	IS: D	PTS:	1	REF:	35		
7.	a. b. c.	when jobs have when jobs have when jobs sizes a when all jobs are	the same differen re not k	e size t sizes nown in advand				
	AN	IS: A	PTS:	1	REF:	35		
8.	a.		t	of free memory	c.	n blocks of allocated memory. External fragmentation Internal fragmentation		
	AN	IS: C	PTS:	1	REF:	36		
9.	me a. b.	e keeps the famory. fixed partition al first-fit memory  IS: B	location	n on	c.	dynamic fit memory allocation best-fit memory allocation		
10.								
	AN	IS: D	PTS:	1	REF:	38		
11.	sch Job J1 J2 J3 J4 Blo B1 B2 B3 B4	neme is used, the jos: 10K 20K 30K 10K  ccks: 30K 15K 50K			run.	s 1-4 and memory blocks. Assuming a first-fit		
	a.	Jl			C.	J3		

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b. J2
                                                  d. J4
    ANS: C
                         PTS: 1
                                              REF: 39|40
12. Consider the following space requirements for jobs 1-4 and memory blocks. Assuming a best-fit
    scheme is used, the job,_____, is placed in the last block.
    Jobs:
    J1 10K
    J2 20K
    J3 30K
    J4 10K
    Blocks:
    B1 30K
    B2 15K
    B350K
    B4 20K
    a. J1
                                                  c. J3
                                                  d. J4
    b. J2
                         PTS: 1
    ANS: B
                                              REF: 40
13. The following algorithm can be described as _____.
     1 Set counter to 1
    2 Do while counter <= number of blocks in memory
      If job_size > memory_size(counter)
        Then counter = counter + 1
      Else
        load job into memory_size(counter)
        adjust free/busy memory lists
        go to step 4
    End do
    3 Put job in waiting queue
    4 Go fetch next job
    a. first-fit memory allocation
                                                  c. least-fit memory allocation
    b. best-fit memory allocation
                                                  d. fixed partition memory allocation
    ANS: A
                         PTS: 1
                                              REF: 41
14. The following algorithm can be described as _____.
     1 Initialize memory_block(0) = 99999
    2 Compute initial_memory_waste = memory_block(0) - job_size
    3 Initialize subscript = 0
    4 Set counter to 1
    5 Do while counter <= number of blocks in memory
         If job_size > memory_size(counter)
            Then counter = counter + 1
         Else
            memory_waste = memory_size(counter) - job_size
         If initial_memory_waste > memory_waste
            Then subscript = counter
            initial memory waste = memory waste
            counter = counter + 1
```

	Else Load job into r		•	least-fit memory allocation fixed partition memory allocation		
	ANS: B	PTS: 1	REF:	42		
15.		eginning address	of the block g Block Size 105 5 600 20 205 4050 230 1000		3	
16.	<ul><li>b. Memory Manag</li><li>c. Memory Manag</li><li>d. Memory Manag</li><li>"free."</li></ul>	ger releases the bl ger immediately g ger adds block to ger resets the statu	ock and comb gives memory free list and re us of the memo	bines it with another free block.  to another program. removes it from busy list. nory block where the job was stored to		
	ANS: D	PTS: 1	REF:	44		
17.	<ul> <li>In a dynamic partition scheme,, is how the Memory Manager deallocates a block that is betwee two other free blocks?</li> <li>a. The sizes of the three free partitions must be combined.</li> <li>b. All three are moved individually from the busy list to the free list.</li> <li>c. The block is combined with the larger of the two adjacent blocks.</li> <li>d. The status of the block is set to free.</li> </ul>					
	ANS: A	PTS: 1	REF:	46		
18.	When memory is de a. blank line b. null entry ANS: B	eallocated, an entr	c.	ioved from the free list by creating a(n) joined entry empty entry		
4.0						
19.	A(n) in the bu returned to the free a. blank line		-	block between two other busy memory blocks is joined entry	}	

## Understanding Operating Systems 5th Edition McHoes Test Bank b. null entry d. empty entry ANS: B PTS: 1 **REF: 47** 20. The \_\_\_\_ of memory, sometimes referred to as garbage collection or defragmentation, is performed by the operating system to reclaim fragmented sections of the memory space. a. deallocation c. compaction b. redirection d. reallocation ANS: C PTS: 1 REF: 48 21. The operating system can tell the \_\_\_\_ of each group of digits by its location in the line and the operation code. a. function c. order b. value d. assignment ANS: A PTS: 1 REF: 49 22. The \_\_\_\_ contains the value that must be added to each address referenced in the program so it will be able to access the correct memory addresses after relocation. a. busy list c. relocation register b. compaction monitor d. bounds register ANS: C PTS: 1 **REF: 52** 23. \_\_\_\_ is the actual memory address for a job that starts at 18K. a. 1,800 c. 18,432 b. 18,000 d. 180,000 ANS: C PTS: 1 **REF: 53** 24. By compacting and relocating, the Memory Manager optimizes the use of memory and thus improves throughput, but an unfortunate side effect is more \_\_\_\_\_. a. null entries c. errors b. segmentation d. overhead ANS: D PTS: 1 **REF: 53**

25. One approach to performing compaction is to do it when a certain of memory becomes busy.

c. bit

REF: 53

d. area

PTS: 1

a. byte

ANS: B

b. percentage